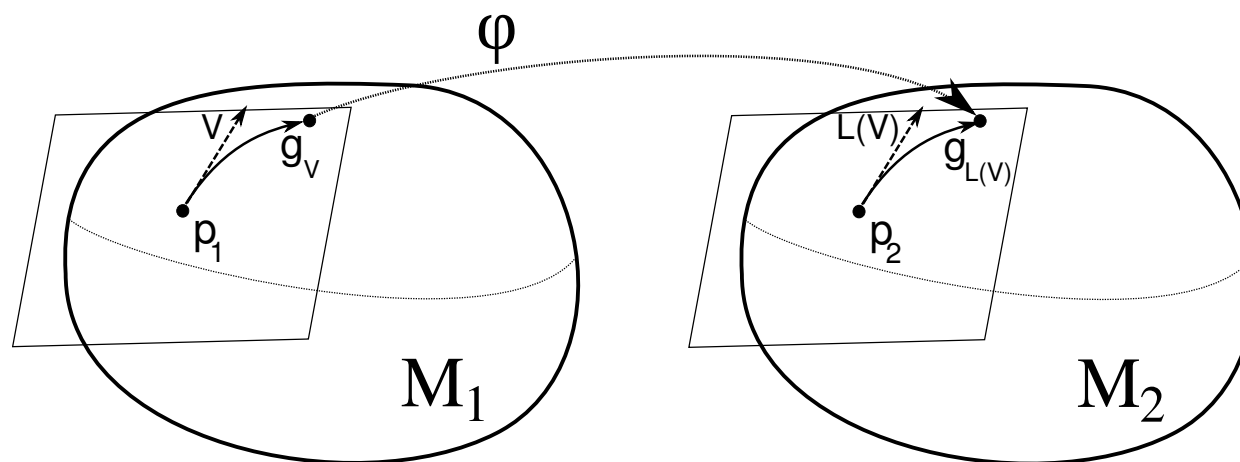


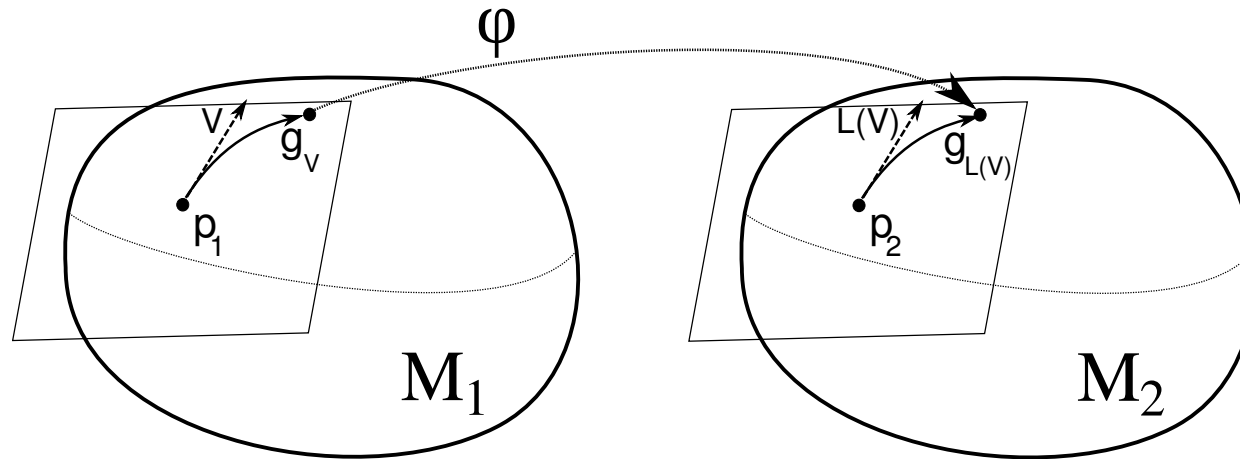
LINKING CURVES, SUTURED MANIFOLDS AND THE AMBROSE CONJECTURE FOR GENERIC 3-MANIFOLDS

Pablo Angulo (ECM7 Berlin, 2016)



The Ambrose problem

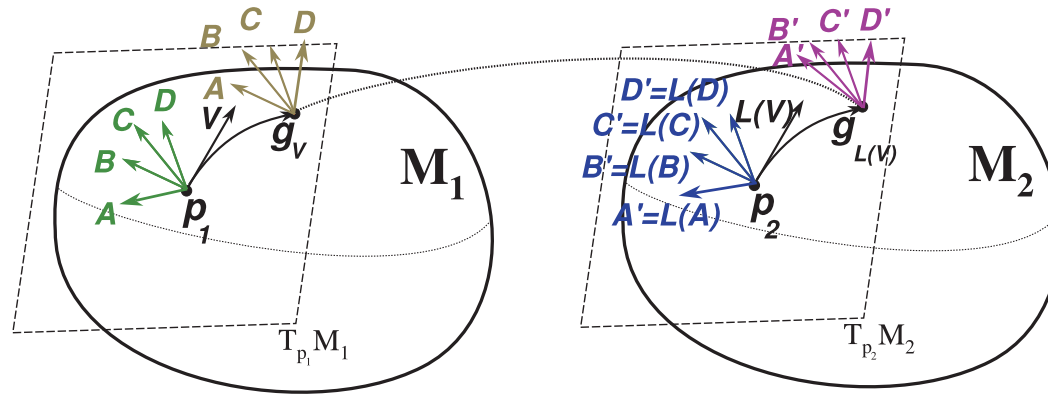
An **isometry** between Riemannian manifolds is determined by its differential at one point p_1 :

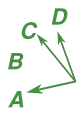








$\varphi: M_1 \rightarrow M_2$ is an isometry between connected and complete manifolds st
 $\varphi(p_1) = p_2, L = d_{p_1}\varphi$

- Fix a *normal neighborhood* U of p_1 .
- Any point q of U can be reached from p_1 by a unique geodesic contained in U , with speed V (parametrized by $[0, 1]$).
- $\varphi(q)$ is the endpoint of the geodesic with initial conditions p_2 and $L(V)$.

Parallel transport of curvature

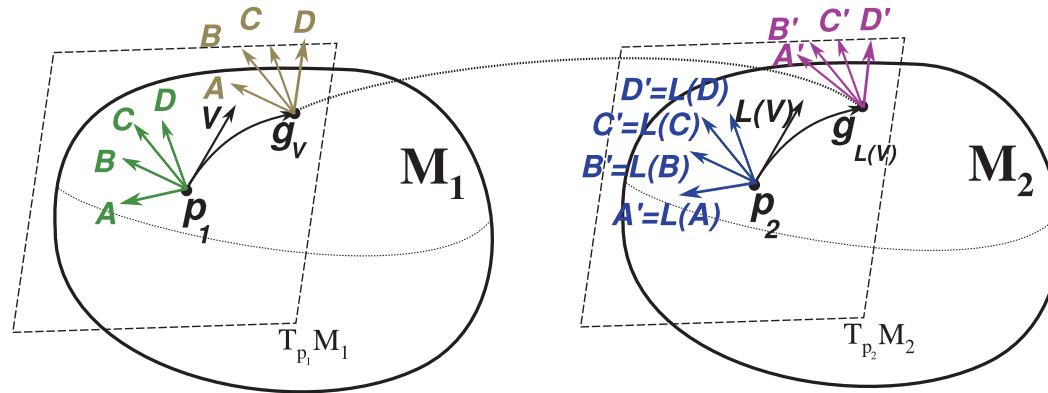


- We parallel displace  along g_V up to $q = g_V(1)$ and compute $R_1(V, \text{) = \langle R_{g_V(1)}(A, B)C, D \rangle$
- We apply L to the vectors: , parallel displace them along $g_{L(V)}$ up to $\varphi(q) = g_{L(V)}(1)$ and compute $R_2(V, \text{) = \langle R_{g_{L(V)}(1)}(A', B')C', D' \rangle$
- If φ is an isometry, $R_1(V, \text{) = R_2(V, \text{)$ for any V and .

Cartan's lemma

Let's do the opposite thing:

let $L: T_{p_1}M_1 \rightarrow T_{p_2}M_2$ be a linear isometry \Rightarrow we get a map φ_L defined in a convex neighborhood of p_1



Cartan's lemma

If $R_1(V, \begin{matrix} B \\ C \\ D \\ A \end{matrix}) = R_2(V, \begin{matrix} B' \\ C' \\ D' \\ A' \end{matrix})$

for any A, B, C, D, V , then φ_L is a local isometry from a convex neighborhood of p_1 to a neighborhood of p_2

Global version of Cartan's lemma

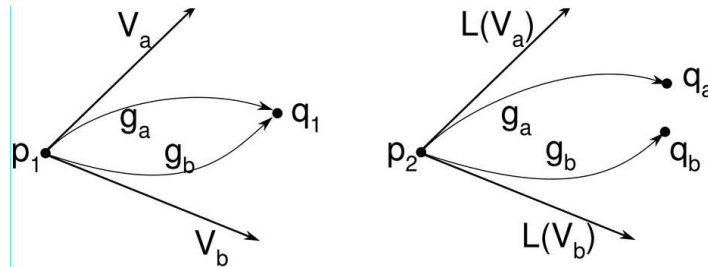
(Cartan-)Ambrose(-Hicks) theorem

If parallel transport of the curvature of M_1 y M_2
along geodesics *with one elbow* coincide,
and both manifolds are *simply connected*,
 φ_L is an isometry from M_1 onto M_2

Ambrose Conjecture (1956)

If parallel transport of the curvature of M_1 y M_2
along *smooth* geodesics coincide,
and both manifolds are *simply connected*,
then φ_L is an isometry from M_1 onto M_2

The proof of the Cartan lemma proves the Ambrose conjecture if there are no conjugate points. The problem is: there may be many geodesics from p_1 to q_1 : are the corresponding points in M_2 the same?



History

- Up to 1987, Cartan's lemma is generalized (Hicks59, Hicks66, O'Neill68, Amici-Casciaro86, BlumenthalHebda87, PawelReckziegel02). The global version (involving geodesics with one elbow), is automatic.
- In **1987, James Hebda proved the conjecture for surfaces**, assuming that the «distance to the cut locus» is an absolutely continuous function.
- J. Hebda (1994) and J-I Itoh (1996) prove independently that this function is indeed absolutely continuous for any smooth surface.
- J-I Itoh and M. Tanaka (2000) prove that it is indeed Lipschitz for a manifold of any dimension, but they could not prove the conjecture from this, since Hebda's proof for surfaces does not work in higher dimensions.
- J. Hebda (2010) proves the conjecture for *generic riemannian manifolds*. The proof does not “pass to the limit”, so new strategies are needed.

Quick review of 1987 James Hebda's proof

- Cartan's lemma provides an isometric immersion of $M_1 \setminus \text{Cut}_{p_1} = \exp_{p_1}(O_{p_1})$ into M_2

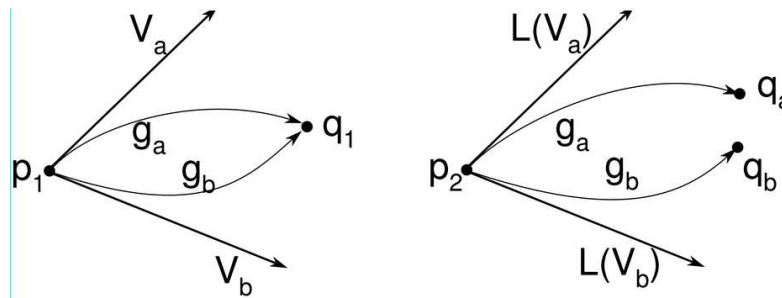
Injectivity set. $O_{p_1} = \{x \in T_{p_1}M_1: d_{M_1}(\exp_{p_1}(p_1, t x)) = t\}$

Tangent cut locus. $\text{TCut}_{p_1} = \partial O_{p_1}$

Cut locus. $\text{Cut}_{p_1} = \exp_{p_1}(\text{TCut}_{p_1})$

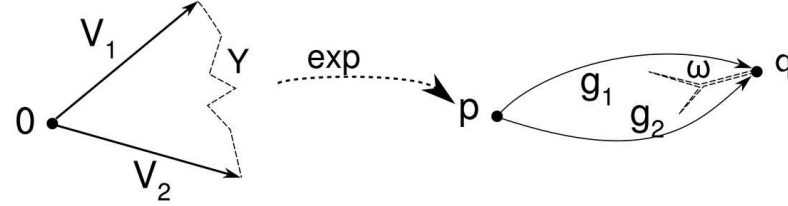
$$\text{Cut}_{p_1} = \left\{ q \in M_1: \begin{array}{l} \text{there are at least two minimizing geodesics from } p_1 \text{ to } q \\ \text{or the unique minimizing geodesic is focal} \end{array} \right\}$$

- If a point $q_1 = \exp_{p_1}(V_a) = \exp_{p_1}(V_b)$ is connected to p_1 by *two minimizing geodesics*, there are two possible images: $\varphi(q_1)$ can be $\exp_{p_2}(L(V_a))$ or $\exp_{p_2}(L(V_b))$. The goal is to prove that it's the same point.



Tree-formed (or tree-like) paths

A central part of the strategy is to find a path Y in TCut_p that joins V_a and V_b (for any point $q = \exp_p(V_a) = \exp_p(V_b) \in \text{Cut}_p$), and maps by \exp_p to a **tree-formed curve** (contained in Cut_p)



Definition 1. Let $u: [0, 1] \rightarrow M$ be an absolutely continuous curve. It is **fully tree formed** if:

- $\exists T: [0, 1] \rightarrow \Gamma$, a quotient map with $T(0) = T(1)$
- u factors through T ($u = \bar{u} \circ T$ for $\bar{u}: \Gamma \rightarrow M$)
- for any continuous 1-form $\varphi(x) \in T_{\bar{u}(x)}^* M$:

$$\int_0^1 \varphi(T(s))(u'(s)) ds = 0$$

Remark: Tree-formed paths reappeared later in the theory of *Rough Path*.

It is easy to find curves in the **tangent cut locus** whose image by \exp is tree-formed in 2D (there are one or two choices, both work)

but in 3D and higher, this is not possible...

Look for linking curves outside the cut locus

$$e_1 = \exp_{p_1}; \quad e_2 = L \circ \exp_{p_2}$$

Definition 2. A ***linking curve*** is an absolutely continuous curve $Y: [0, l] \rightarrow T_p M$ such that $e_1 \circ Y$ is a fully tree formed curve.

Definition 3. Two points $x, y \in T_p M$ are ***strongly linked*** iff \exists linking curve Y , $Y(0) = x$, $Y(1) = y$

$$x, y \text{ strongly linked} \Rightarrow e_1(x) = e_1(y), e_2(x) = e_2(y)$$

Definition 4. $O \in T_p M$ is ***unequivocal*** iff $e_1(O)$ is open, and there is an isometry $\varphi_O: e_1(O) \rightarrow e_2(O)$ such that $\varphi_O \circ e_1|_O = e_2|_O$

Definition 5. $x \in T_p M$ is ***unequivocal*** if there is a sequence of sets W_n such that $e_1(W_n)$ is a neighborhood basis of $e_1(x)$.

The idea is to link singularities of \exp_p to unequivocal points

Main theorem

Theorem 6. *Let M_1, M_2 be simply connected Riemannian manifolds with L -related curvature, such that **every $x \in V_1$ is linked to some unequivocal $y \in V_1$, with $|y| \leq |x|$.***

*Then there is a **strong synthesis** of M_1 and M_2 through e_1 and e_2 .*

Proof. A manifold M is defined as a **quotient** of a subset of $T_{p_1}M_1$, **identifying linked points**. The maps e_1 and e_2 induce maps π_1 and π_2 .

A topology is defined ad-hoc, we prove that π_1, π_2 are local homeomorphisms, etc.

The condition $|y| \leq |x|$ is important to prove they are covering maps.

Thus π_1 and π_2 are covering maps and local isometries. We say M is a *synthesis* of M_1 and M_2 .

M_1, M_2 simply connected $\Rightarrow \pi_1$ and π_2 are global isometries. □

In other words, *if we can find linking curves* starting at conjugate points,
the Ambrose conjecture follows.

But can we find enough linking curves?

Generic 3D riemannian manifolds

For a *generic* set of riemannian metrics in a 3D manifold, $T_p M$ admits the following decomposition ([canonical form of \$\exp_p\$](#)):

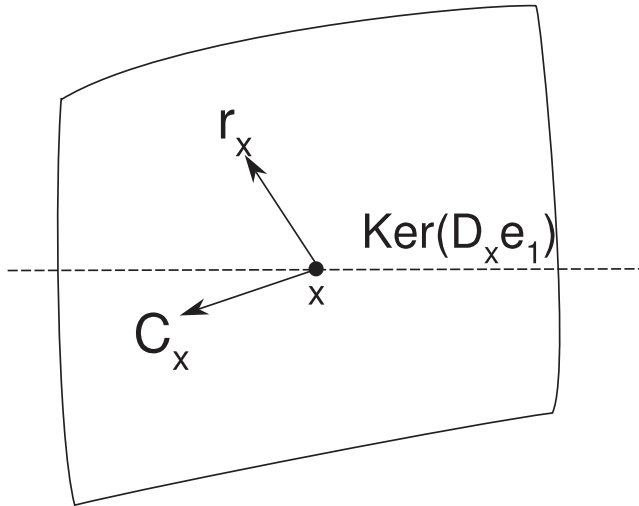
- An open set consisting of non-conjugate points (**NC**) $((x_1, \dots, x_n) \rightarrow (x_1, \dots, x_n))$.
- Strata of dimension **2**, of points of type **A₂** (*fold* singularities)
 $(x_1, x_2, \dots, x_n) \rightarrow (x_1^2, x_2, \dots, x_n)$.
- Strata of dimension **1**, of points of type **A₃** (*cusp* singularities). We further split them into **A₃(I)** and **A₃(II)** (minima and maxima, roughly)
 $(x_1, x_2, \dots, x_n) \rightarrow (x_1^3 \pm x_1 x_2, x_2, \dots, x_n)$
- Isolated points of type **A₄** $((x_1, x_2, \dots, x_n) \rightarrow (x_1^4 + x_1^2 x_2 + x_1 x_3, x_2, \dots, x_n))$
- Isolated points of type **D₄⁺** $((x_1, x_2, \dots, x_n) \rightarrow (\frac{1}{2}x_1^2 + x_2 x_3, \frac{1}{2}x_2^2 + x_1 x_3, \dots, x_n))$
- Isolated points of type **D₄⁻** $((x_1, x_2, \dots, x_n) \rightarrow (\frac{1}{2}x_1^2 - \frac{1}{2}x_2^2 + x_1 x_3, -x_1 x_2 + x_2 x_3, \dots, x_n))$

Conjugate flow at points of type A_2

At points of type A_2 the kernel of the exponential is transversal to the set of conjugate points (which is a smooth hypersurface). The exponential e_1 is given in adapted coordinates by:

$$(x_1, \dots, x_{n-1}, x_n) \rightarrow (x_1, \dots, x_{n-1}, x_n^2)$$

e_1 is not a local homeomorphism at an A_2 point, so how can we link the A_2 point x to an unequivocal point y ? The conjugate flow tells us how we can start:



Gauss' lemma implies that the radial vector r_x is transversal to the kernel of the exponential. The sum of both spaces is a plane that we intersect with the tangent to the set of conjugate points:

$$(\ker D_x e_1 \oplus \langle r_x \rangle) \cap T \text{Conj} = \langle C_x \rangle$$

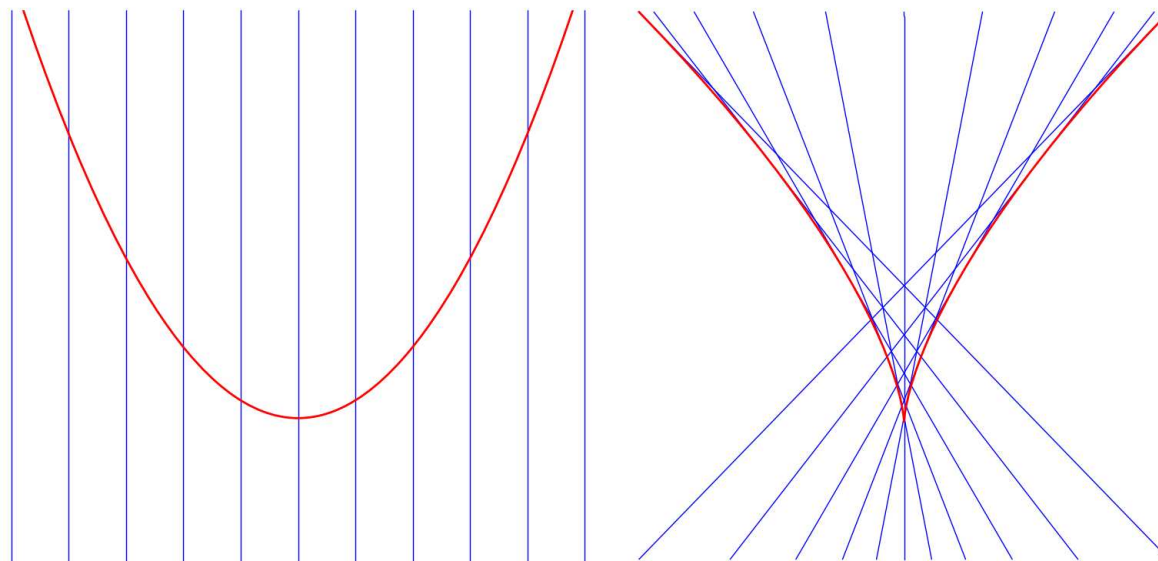
We choose C_x such that $C_x \cdot r_x < 0$

Conjugate flow

A **conjugate descending curve (CDC)** is an *integral curve of the vector field C_x* . The curve stays within the set of conjugate points, and can be continued until it hits a point that is not A_2 .

The most simple and most important case is that the conjugate flow hits an **A_3** point. For A_3 points, e_1 is given in adapted coordinates by:

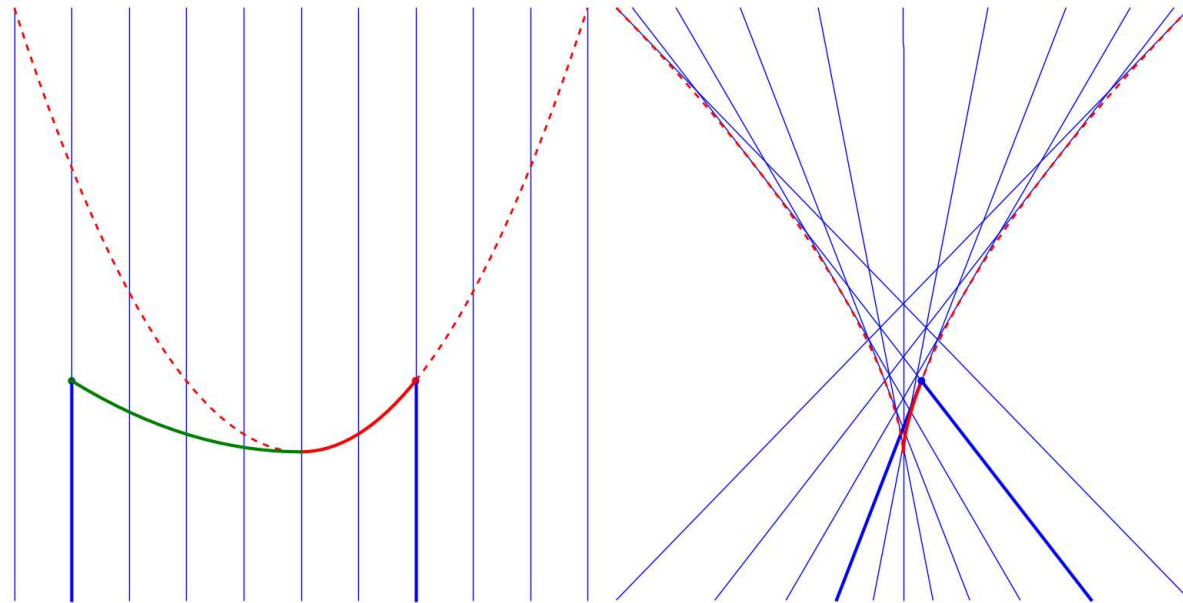
$$(x_1, \dots, x_{n-1}, x_n) \rightarrow (x_1, \dots, x_{n-1}, x_n^3 + x_1 x_n)$$



Replying a CDC

Once the CDC reaches an A_3 point, we can find a **reply** to the CDC: it is a **curve of NC points** whose image by e_1 is the same as the CDC, but run in the opposite direction.

The concatenation of both segments is a curve **whose image is tree-formed**.



A CDC α is **unbeatable**: suppose β replies to α :

$$|\alpha(0)| - |\alpha(t_0)| = \text{length}(\exp \circ \alpha) = \text{length}(\exp \circ \beta) > |\beta(t_0)| - |\beta(0)|$$

Building linking curves

Remember: *our goal is to show our A_2 point x is linked to an unequivocal point y*

Starting at x , suppose we follow the CDC up to an A_3 point. We can keep replying as long as the reply stays within V_1 , but **we may hit a singularity**. If that happens, we descend along the conjugate descending flow again, etcetera:

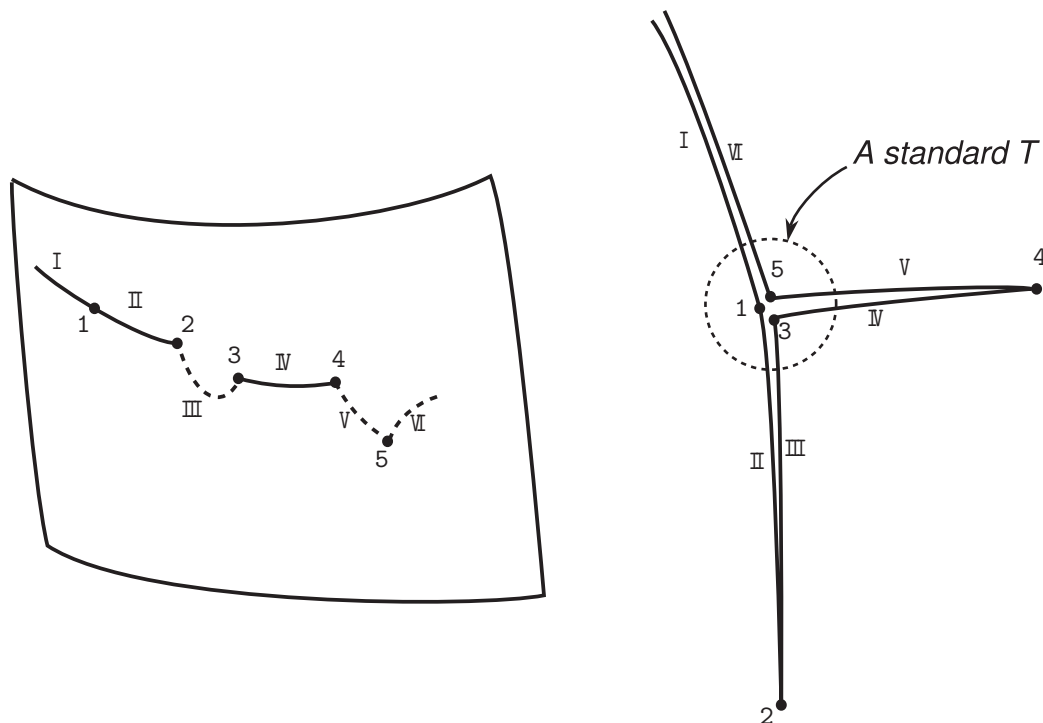
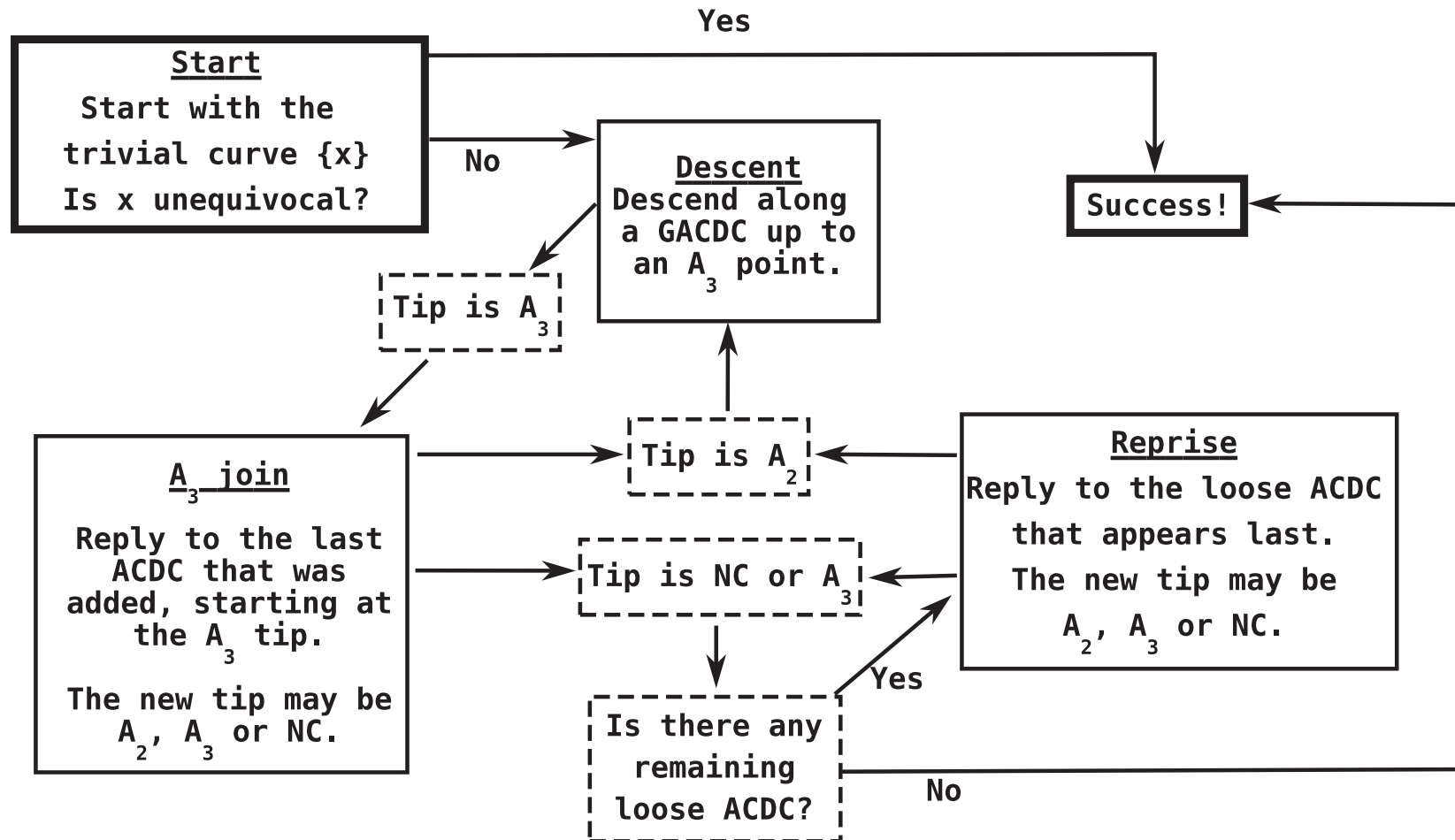


Figure. Splitter (1), A_3 -join (2,4), hit(3), reprise(5)

The algorithm

We summarize the procedure in this flow diagram:



Does it work?

- *Do we always reach an A_3 point?* Yes, we can «dodge» other singularities.
- *How long can we keep replying?* We can always reply to an A_2 point, since the «linking curve under construction» always stays within V_1 .
- *What if the reply hits a worse singularity?* We can also dodge that.
- *How do you know the procedure will ever stop?* Because every point in a generic manifold has a «**transient neighborhood**»: if the algorithm starts at a point on that neighborhood, after a finite number of elementary steps of the algorithm, the tip will never again be there (the *unbeatable* property is key).

We show that it works for generic metrics, in 3 dimensions. The algorithm can also start at $A_3(\text{II})$, A_4 , D_4^+ and D_4^- points.

Points of type $A_3(\text{I})$ are unequivocal.

Summary

- For a generic manifold, $V_1 \subset \text{NC} \cup A_2 \cup A_3(I) \cup A_3(\text{II}) \cup A_4 \cup D_4^\pm$
- $\mathcal{I} = V_1 \cap (\text{NC} \cup A_3(I)) \quad \mathcal{J} = V_1 \cap (A_2 \cup A_3(\text{II}) \cup A_4 \cup D_4^\pm)$

Theorem 7. *A point in \mathcal{I} is unequivocal; a point in \mathcal{J} is linked to a point of \mathcal{I} of smaller radius.*

Corollary 8. *Ambrose conjecture holds for a generic 3-manifold.*

Thanks for your attention!

- Pablo Angulo. Linking curves, sutured manifolds and the Ambrose conjecture for generic 3-manifolds. arxiv.org/abs/1509.02125

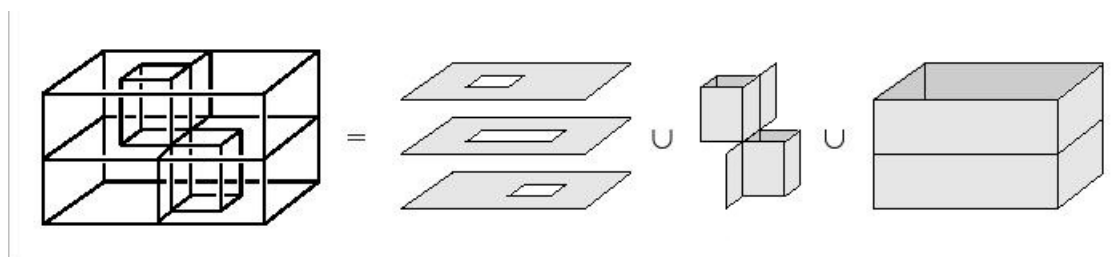
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FAQ #1 Why can't you extend the proof of James Hebda to three dimensions?

Answer: You mean: Can we find the curve Y in dimension ≥ 3 ?

In **dimension ≥ 3** there are many possible choices for Y , none of them is *canonical*.

But it's worse than that: the house with two rooms is the cut locus of a certain manifold, and it doesn't have edge points!

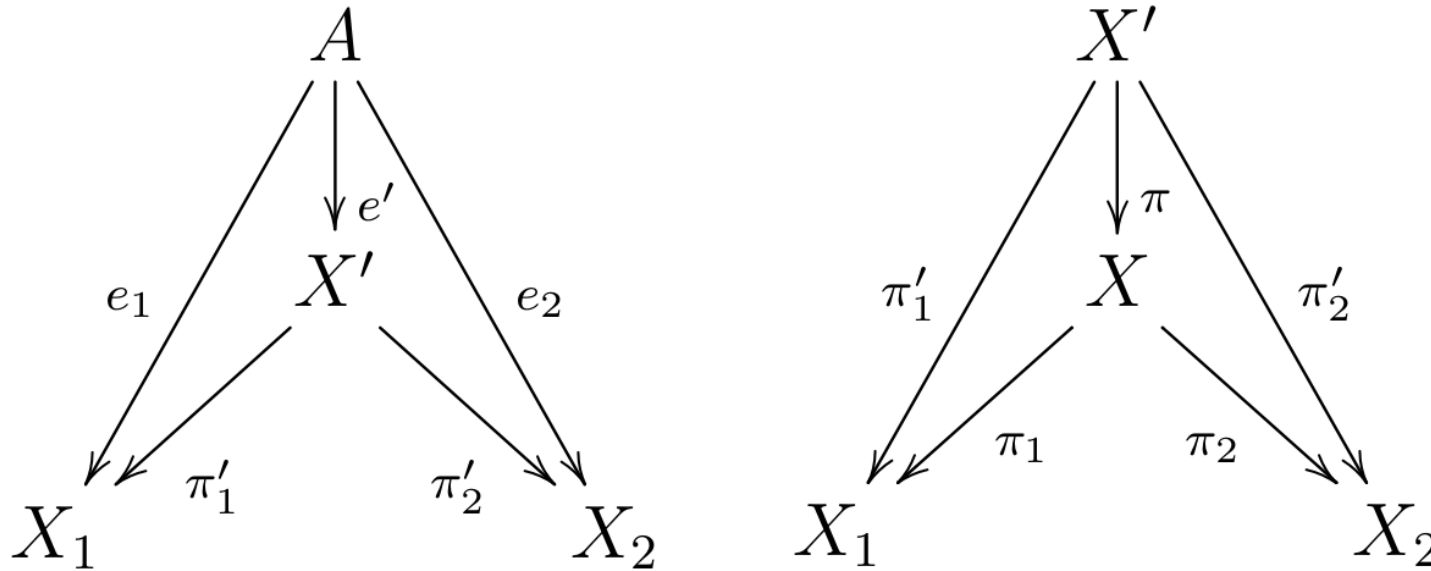


Weinstein showed that every manifold except S^2 admits a metric such that the cut locus wrta point doesn't have edge points.

There cannot be non-trivial linking curves contained in the cut locus of such manifolds.

FAQ #2 Is the synthesis the «minimal common covering space»?

Yes, it satisfies a universal property (5.4.7): Let X be the synthesis of X_1 and X_2 . For any Riemannian manifold X' , continuous surjective map $e': A \rightarrow X'$ and local isometries $\pi'_1: X' \rightarrow X_1$ and $\pi'_2: X' \rightarrow X_2$, such that $e_i = \pi'_i \circ e'$, for $i = 1, 2$, there is a local isometry $\pi: X' \rightarrow X$ such that:



FAQ #3 What is the motivation for studying Ambrose Problem?

Ambrose's motivation (roughly): to characterize a Riemannian manifold by the parallel transport of its curvature: $L: \mathbb{R}^n \times G_2(\mathbb{R}^n) \rightarrow \mathbb{R}$

$L(V, P)$: parallel transport the plane P along the geodesic g_V and compute the sectional curvature of the plane.

My motivation: The Ambrose problem is similar to some *inverse problems* in Riemannian geometry. Some of these problems come directly from applications like tomography. Some of those problems are easy to solve if there are no singularities. I planned to build some muscle and try other such inverse problems later.

FAQ #4 Can you bound the lenght of the linking curves and approximate an arbitrary metric by generic ones?

Not without new ideas. I've tried two things:

1. As the *slack* of an A_2 point (the angle between the kernel of the exponential and the tangent hyperplane to the conjugate points) goes to 0, and x becomes a worse singularity, the *gain* of the CDC α through x decreases. It's true that the lenght of its composition with the exponential also decreases, but overall, nor the lenght of α neither that of $e_1 \circ \alpha$ is bounded.
2. Try to put numbers to the algorithm: Let B_R be the maximum lenght of a linking curve through a point x of radius R . The algorithm starts with a CDC α of lenght l that leaves a transient neighborhood U of x . Then a linking curve starting at the tip of α follows. Its lenght is bounded by B_{R-l} . And after that, we have to reply to α . If we can reply to α at once:

$$B_R < l + B_{R-l}$$

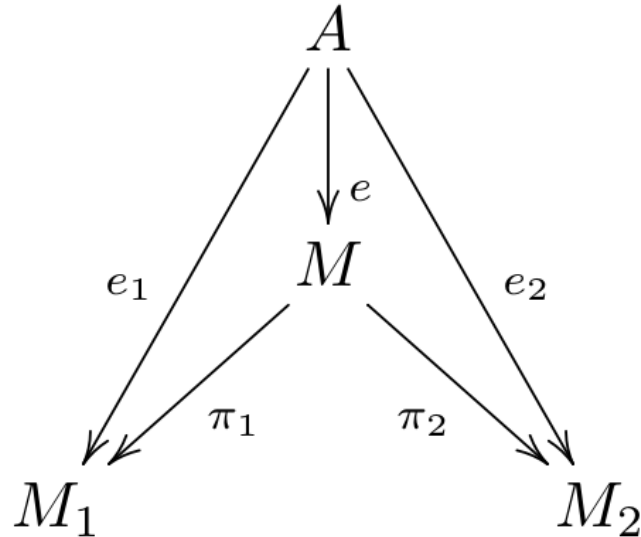
but the reply might hit an A_2 point, and then we have to reply to part of α , then plug in another linking curve at the tip of the reply, then reply the rest... if there are k interruptions:

$$B_R < 2l + k B_{R-\varepsilon}$$

where ε is the *gain* of the transient neighborhood. This is exponential growth.

FAQ #5 Why “*strong* synthesis”? Is there a “*weak synthesis*”?

Yes, if π_1 and π_2 are local isometries, but not covering maps:



The corresponding theorem is:

Theorem 9. *Let M_1, M_2 be Riemannian manifolds with L -related curvature, such that every $x \in V_1$ is linked to some unequivocal $y \in V_1$.*

*Then there is a **strong synthesis** of M_1 and M_2 through e_1 and e_2 .*

In short: if we do not assume we can find y with $|y| \leq |x|$, then π_1 and π_2 may not be covering maps.

FAQ #6 Why “*strongly* linked”? Are there “*weakly linked*” points?

Yes, but that is rather technical. Let me just say that there are other hypothesis that imply:

$$e_1(x) = e_1(y), e_2(x) = e_2(y)$$